

Sorting Algorithms

1. Seleksi

2. Gelembung

3. Penyisipan

4. Merge

5. Quick

Metode Seleksi

(one of the simplest sorting algorithms)

3	10	4	6	8	9	7	2	1	5
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Cari dalam keseluruhan array, temukan nilai terbesar, (10) dan tukar nilai ini dengan nilai yang tersimpan dalam lokasi terakhir dari array (5)

3	10	4	6	8	9	7	2	1	5
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3	5	4	6	8	9	7	2	1	10
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Temukan nilai terbesar kedua dalam array (9), tukar dengan nilai yang tersimpan dalam lokasi terakhir kedua(1).

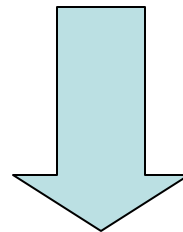
3	5	4	6	8	9	7	2	1	10
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3	5	4	6	8	1	7	2	9	10
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3	5	4	6	8	1	7	2	9	10
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Dua nilai terakhir yang bertanda biru merupakan posisi yang pasti karena keduanya merupakan nilai terbesar dan nilai terbesar kedua.

Sekarang, ulangi proses “seleksi dan tukar” ...



3	5	4	6	8	1	7	2	9	10
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3	5	4	6	2	1	7	8	9	10
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3	5	4	6	2	1	7	8	9	10
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3	5	4	6	2	1	7	8	9	10
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3	5	4	1	2	6	7	8	9	10
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3	5	4	1	2	6	7	8	9	10
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3	2	4	1	5	6	7	8	9	10
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3	2	4	1	5	6	7	8	9	10
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3	2	1	4	5	6	7	8	9	10
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3	2	1	4	5	6	7	8	9	10
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1	2	3	4	5	6	7	8	9	10
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1	2	3	4	5	6	7	8	9	10
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Algoritma

1. Ulangi langkah 2-7 untuk $i=0$ s/d $n-2$
2. $\text{min} = \text{data}[i]$
3. Ulangi langkah 4 untuk $j=i+1$ s/d $n-1$
4. Lakukan langkah 5-6 jika $\text{data}[j] < \text{min}$
5. $\text{min} = \text{data}[j]$
6. $\text{min_id} = j$
7. tukar $\text{data}[i]$ dengan $\text{data}[\text{min_id}]$

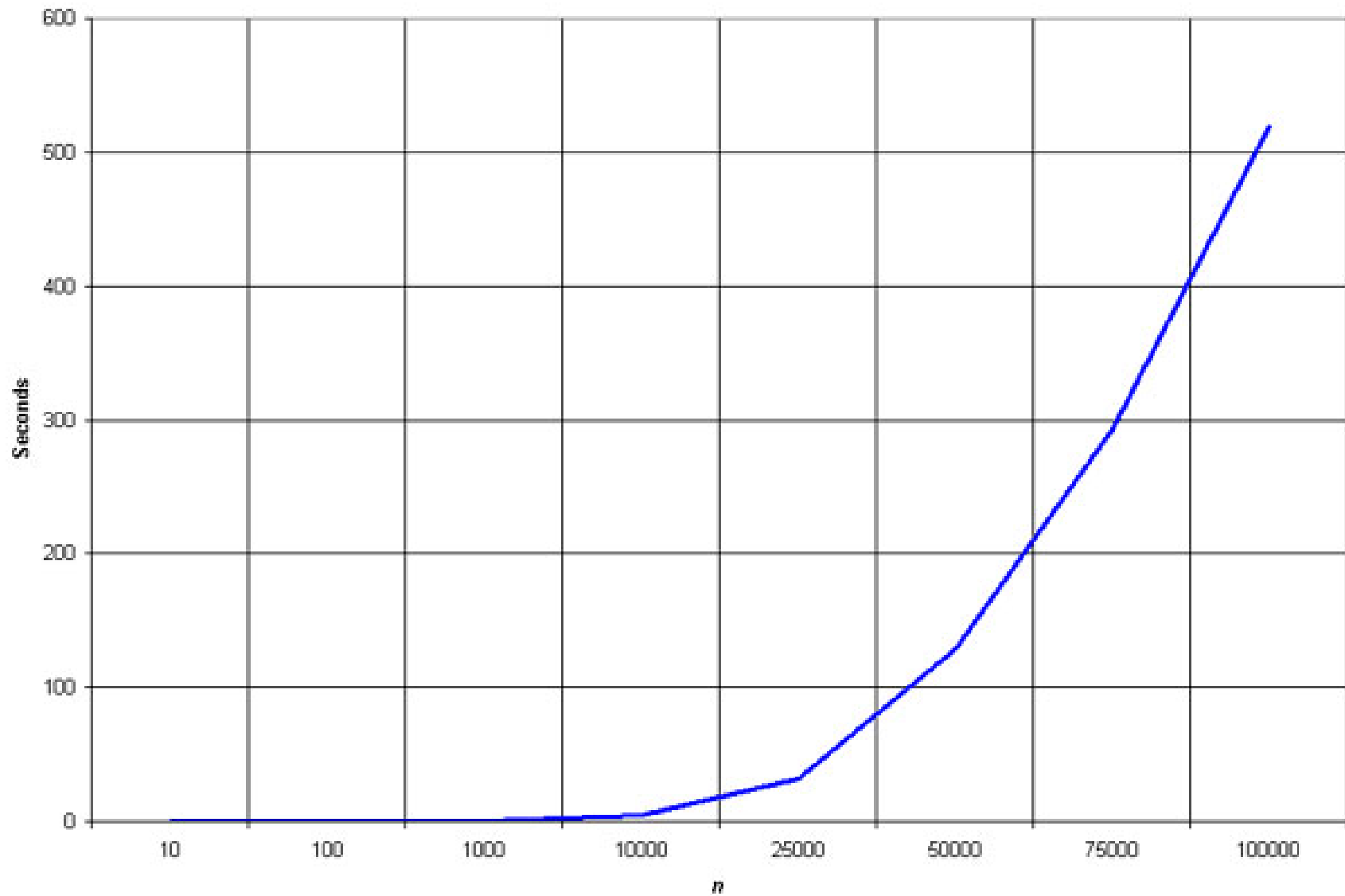
Selection Sort → Analysis

- Moves from left to right, putting elements into their final position without looking back
- Wastes most of the time looking for the minimum element in the unsorted part of the array

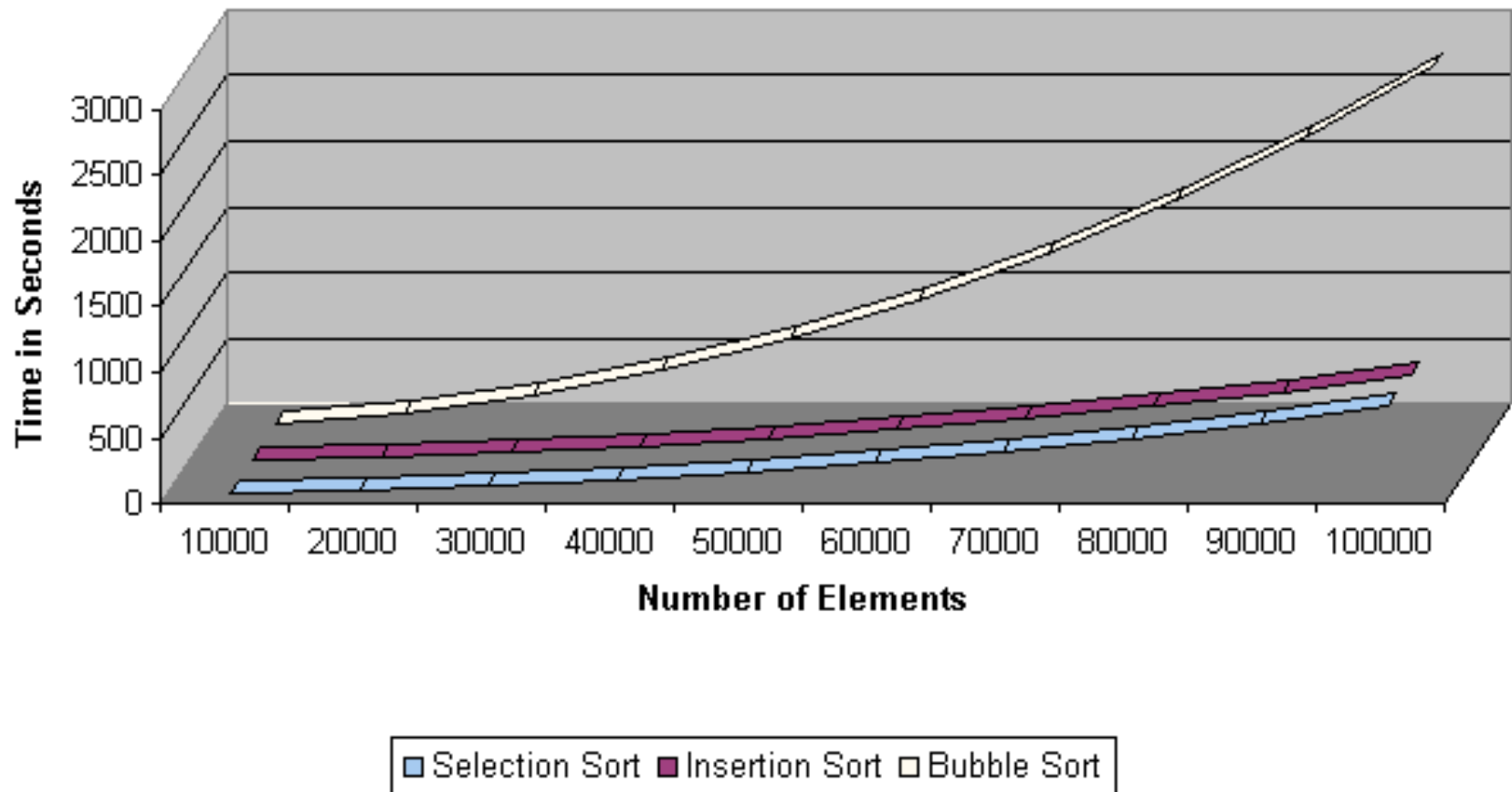
Selection Sort → Analysis

- $O(n^2)$ for worst and average cases
- Uses about $n^2/2$ comparisons
- Uses about n exchanges

Selection Sort → Empirical Analysis



Running Times of Sorting Algorithms



Selection Sort yields a 60% performance improvement over the Bubble Sort